2.2 Heuristic Instantiation

DPLL(T) is limited to ground (or existentially quantified) formulas. Even if we have decidability for more than the ground fragment of a theory $\mathcal{T}$, we cannot use this in DPLL(T).

Most current SMT implementations offer a limited support for universally quantified formulas by heuristic instantiation.

Goal:

Create potentially useful ground instances of universally quantified clauses and add them to the given ground clauses.

Idea (Detlefs, Nelson, Saxe: Simplify):

Select subset of the terms (or atoms) in $\forall\vec{x}\,C$ as "trigger" (automatically, but can be overridden manually).

If there is a ground instance $C\theta$ of $\forall\vec{x}\,C$ such that $t\theta$ occurs (modulo congruence) in the current set of ground clauses for every $t \in \mathit{trigger}(C)$, add $C\theta$ to the set of ground clauses (incrementally).

Conditions for trigger terms (or atoms):

(1) Every quantified variable of the clause occurs in some trigger term (therefore more than one trigger term may be necessary).

(2) A trigger term is not a variable itself.

(3) A trigger is not explicitly forbidden by the user.

(4) There is no larger instance of the term in the formula:
(If $f(x)$ were selected as a trigger in $\forall x\, P(f(x), f(g(x)))$, a ground term $f(a)$ would produce an instance $P(f(a), f(g(a)))$, which would produce an instance $P(f(g(a)), f(g(g(a))))$, and so on.)

(5) No proper subterm satisfies (1)–(4).

Also possible (but expensive, therefore only in restricted form): Theory matching

The ground atom $P(a)$ is not an instance of the trigger atom $P(x+1)$; it is however equivalent (in linear algebra) to $P((a-1)+1)$, which *is* an instance and may therefore produce a new ground clause.

Heuristic instantiation is obviuosly incomplete

e. g., it does not find the contradiction for $f(x,a) \approx x$, $f(b,y) \approx b$, $a \not\approx b$

but it is quite useful in practice:

modern implementations: CVC, Yices, Z3.

2.3 Local Theory Extensions

Under certain circumstances, instantiating universally quantified variables with "known" ground terms is sufficient for completeness.

Scenario:

$\Sigma_0 = (\Omega_0, \Pi_0)$: base signature;
$\mathcal{T}_0$: Σ_0-theory.

$\Sigma_1 = (\Omega_0 \cup \Omega_1, \Pi_0)$: signature extension;
K: universally quantified Σ_1-clauses;
G: ground clauses.

Assumption: clauses in G are Σ_1-flat and Σ_1-linear:

only constants as arguments of Ω_1-symbols,

if a constant occurs in two terms below a Ω_1-symbol, then the two terms are identical,

no term contains the same constant twice below below a Ω_1-symbol.

Example: Monotonic functions over $\mathbb{Z}$.

$\mathcal{T}_0$: Linear integer arithmetic.

$\Omega_1 = \{f/1\}$.
$K = \{\ \forall x, y\ (\neg x \leq y\ \vee\ f(x) \leq f(y))\ \}$.

$G = \{\ f(3) \geq 6,\ f(5) \leq 9\ \}$.

Observation: If we choose interpretations for $f(3)$ and $f(5)$ that satisfy the G and monotonicity axiom, then it is always possible to define f for all remaining integers such that the monotonicity axiom is satisfied.

Example: Strictly monotonic functions over $\mathbb{Z}$.

$\mathcal{T}_0$: Linear integer arithmetic.

$\Omega_1 = \{f/1\}$.
$K = \{\ \forall x, y\ (\neg x < y\ \vee\ f(x) < f(y))\ \}$.

$G = \{\ f(3) > 6,\ f(5) < 9\ \}$.

Observation: Even though we can choose interpretations for $f(3)$ and $f(5)$ that satisfy G and the strict monotonicity axiom (map $f(3)$ to 7 and $f(5)$ to 8), we cannot define $f(4)$ such that the strict monotonicity axiom is satisfied.

To formalize the idea, we need partial algebras:

like (usual) total algebras, but $f_{\mathcal{A}}$ may be a partial function.

There are several ways to define equality in partial algebras (strong equality, Evans equality, weak equality, etc.). Here we use weak equality:

an equation $s \approx t$ holds w. r. t. $\mathcal{A}$ and β if both $\mathcal{A}(\beta)(s)$ and $\mathcal{A}(\beta)(t)$ are defined and equal or if at least one of them is undefined;

a negated equation $s \not\approx t$ holds w. r. t. $\mathcal{A}$ and β if both $\mathcal{A}(\beta)(s)$ and $\mathcal{A}(\beta)(t)$ are defined and different or if at least one of them is undefined.

If a partial algebra $\mathcal{A}$ satisfies a set of formulas N w. r. t. weak equality, it is called a weak partial model of N.

A partial algebra $\mathcal{A}$ embeds weakly into a partial algebra $\mathcal{B}$ if there is an injective total mapping $h : U_{\mathcal{A}} \to U_{\mathcal{B}}$ such that if $f_{\mathcal{A}}(a_1, \dots, a_n)$ is defined in $\mathcal{A}$ then $f_{\mathcal{B}}(h(a_1), \dots, h(a_n))$ is defined in $\mathcal{B}$ and equal to $h(f_{\mathcal{A}}(a_1, \dots, a_n))$.

A theory extension $\mathcal{T}_0 \subseteq \mathcal{T}_0 \cup K$ is called *local*, if for every set G, $\mathcal{T}_0 \cup K \cup G$ is satisfiable if and only if $\mathcal{T}_0 \cup K[G] \cup G$ has no (partial) model, where $K[G]$ is the set of instances of clauses in K in which all terms starting with an Ω_1-symbol are ground terms occurring in K or G.

If every weak partial model of $\mathcal{T}_0 \cup K$ can be embedded into a a total model, then the theory extension $\mathcal{T}_0 \subseteq \mathcal{T}_0 \cup K$ is local (Sofronie-Stokkermans 2005).

Note: There are many variants of partial models and embeddings corresponding to different kinds of locality.

Examples of local theory extensions:

free functions

constructors/selectors

monotonic functions

Lipschitz functions.