

**Assignment 1 (CDCL)**

(8 + 4 + 4 = 16 points)

Let  $N$  be the following set of propositional clauses:

$$P \qquad \qquad \qquad \vee \qquad S \qquad \qquad \qquad (1)$$

$$P \vee Q \qquad \qquad \qquad \vee \neg T \qquad \qquad \qquad (2)$$

$$Q \qquad \vee \neg S \qquad \qquad \vee \neg U \qquad \qquad \qquad (3)$$

$$\neg Q \vee R \vee S \vee \neg T \vee U \qquad \qquad \qquad (4)$$

$$R \vee \neg S \qquad \qquad \qquad (5)$$

$$R \qquad \qquad \vee T \qquad \qquad \qquad (6)$$

$$\neg P \qquad \qquad \vee S \qquad \qquad \vee \neg U \qquad \qquad \qquad (7)$$

$$\neg R \vee \neg S \qquad \qquad \qquad (8)$$

$$\neg R \qquad \qquad \vee T \vee U \qquad \qquad \qquad (9)$$

$$\neg P \qquad \vee \neg R \vee S \vee \neg T \qquad \qquad \qquad (10)$$

**Part (a)**Use the CDCL procedure to compute a model of  $N$ .**Part (b)**Use the final state of the CDCL procedure in part (a) to determine the number of models of  $N$  in which  $P$  is true. Give a brief explanation.**Part (c)**Use the CDCL procedure to show that  $N$  has no model in which  $P$  is false.

For both part (a) and (c):

Use the CDCL inference rules with a reasonable strategy. If you use the *Decide* rule, use the largest undefined positive literal according to the ordering  $P > Q > R > S > T > U$ . If you use the *Backjump* rule, determine a suitable backjump clause using the 1UIP method and use the best possible successor state for that backjump clause.

**Assignment 2 (Resolution)**

(8 + 8 = 16 points)

Let  $\Sigma = (\Omega, \Pi)$  be a signature, let  $N$  be a set of ground  $\Sigma$ -clauses.**Part (a)**Prove or refute: All clauses in  $Res^*(N)$  are tautologies if and only if all clauses in  $N$  are tautologies.**Part (b)**Prove or refute:  $Res(N)$  is satisfiable if and only if  $N$  is satisfiable.

**Assignment 3** (*Propositional Logic*) (12 points)

Let  $\Pi$  be a set of propositional variables. Let  $N$  and  $N'$  be sets of clauses over  $\Pi$ . Let  $S$  be a set of literals that does not contain any complementary literals. Prove: If every clause in  $N$  contains at least one literal  $L$  with  $L \in S$  and if no clause in  $N'$  contains a literal  $L$  with  $\bar{L} \in S$ , then  $N \cup N'$  is satisfiable if and only if  $N'$  is satisfiable.

**Assignment 4** (*First-order Logic, Semantics*) (12 points)

Let  $\Sigma = (\Omega, \Pi)$  be a first-order signature with  $\Omega = \{b/0, f/1\}$  and  $\Pi = \{P/1\}$ . Decide for each of the following statements whether they are true or false:

- (1) There is a  $\Sigma$ -model  $\mathcal{A}$  of  $P(b) \wedge \neg P(f(b))$  such that  $U_{\mathcal{A}} = \{7, 8, 9\}$ .
- (2) There is a  $\Sigma$ -model  $\mathcal{A}$  of  $P(b) \wedge \neg P(f(f(b)))$  such that  $f_{\mathcal{A}}(a) = a$  for every  $a \in U_{\mathcal{A}}$ .
- (3)  $P(b) \wedge \neg P(f(b))$  has a Herbrand model.
- (4)  $P(b) \wedge \forall x \neg P(x)$  has a Herbrand model.
- (5)  $\forall x P(f(x))$  has a Herbrand model with a two-element universe.
- (6)  $\forall x P(x)$  has exactly one Herbrand model.
- (7)  $\forall x P(f(x))$  entails  $\forall x P(f(f(x)))$ .

(Note on grading: A yes/no answer is sufficient; you do not have to give an explanation. However, you need at least four correct answers to get any points for this assignment. Missing answers count like false answers.)

**Assignment 5** (*First-order Logic, CNF Transformation*) (12 points)

Let  $\Sigma = (\{c/0, f/1\}, \{P/4, Q/4\})$ . Transform the  $\Sigma$ -formula

$$F = \exists v \forall x \forall y \forall z \neg \forall w \left( \neg P(c, w, z, x) \wedge Q(w, y, f(x), v) \right)$$

into clause normal form using the improved algorithm from Section 3.6. (There are no subformulas in  $F$  for which one should introduce a definition.)

**Assignment 6** (*Clause Orderings*) (12 points)

Find a strict total ordering  $\succ$  on the ground atoms  $P(b)$ ,  $P(c)$ ,  $Q$ ,  $R$  such that

$$P(b) \vee \neg P(c) \succ_C \neg P(b) \vee P(c) \quad (1)$$

$$P(b) \vee P(b) \vee P(b) \vee R \succ_C P(b) \vee R \vee R \quad (2)$$

$$\neg P(b) \vee Q \succ_C P(c) \vee R \quad (3)$$