

Assignment 1 (CDCL)

(14 points)

A Horn clause is a clause with at most one positive literal. Let N be a finite set of propositional Horn clauses and let $\varepsilon \parallel N \Rightarrow_{\text{CDCL}}^* M \parallel N$ be a CDCL derivation in which only the “Unit Propagate” rule is used. Prove: If neither “Unit Propagate” nor “Fail” can be applied to $M \parallel N$, then N is satisfiable.

Assignment 2 (Ordered resolution)

(8 + 4 + 4 = 16 points)

Let $\Sigma = (\{f/1, g/2, b/0, c/0\}, \{P/2, Q/1, R/2\})$; let the atom ordering \succ be a Knuth-Bendix ordering with weight 1 for all function and predicate symbols and variables and the precedence $P > Q > R > f > g > b > c$; let N be the following set of clauses over Σ :

$$\neg R(b, x) \quad (1)$$

$$Q(g(x, z)) \vee Q(g(b, f(y))) \vee \neg R(x, z) \quad (2)$$

$$P(f(x), x) \vee \neg Q(f(c)) \vee \neg R(x, c) \quad (3)$$

$$\neg P(x, c) \vee Q(x) \quad (4)$$

$$\neg P(x, b) \vee R(c, y) \quad (5)$$

Part (a) Suppose that the selection function sel selects no literals. Compute all $\text{Res}_{sel}^>$ inferences between the clauses (1)–(5). (Do not compute inferences between derived clauses. Do not compute any inferences that violate the restrictions of the calculus.)

Part (b) Is the set N saturated up to redundancy? Give a brief explanation.

Part (c) Does a selection function sel exist for which there are no $\text{Res}_{sel}^>$ inferences between the clauses (1)–(5)? If yes, which literals does sel select in the clauses (1)–(5)?

Assignment 3 (E-Algebras)

(4 + 6 = 10 points)

Let $\Sigma = (\Omega, \emptyset)$ with $\Omega = \{f/1, b/0, c/0, d/0\}$; let E be the set of (implicitly universally quantified) equations $\{f(f(x)) \approx b\}$.

Part (a) Show that $b \leftrightarrow_E^* f(b)$. How does the rewrite proof look?

Part (b) Is the universe of the initial E -algebra $\text{T}_\Sigma(\emptyset)/E$ finite or infinite? If it is finite, how many elements does it have?

Assignment 4 (LPO)

(14 points)

Let $\Sigma = (\Omega, \emptyset)$ be a finite signature, let h be a unary function symbol in Ω , and let \succ be a strict partial ordering (“precedence”) on Ω such that h is the smallest element of Ω w.r.t. \succ .

Prove: For all terms $s, t \in T_{\Sigma}(X)$, we have $s \succ_{\text{lpo}} t$ if and only if $s \succeq_{\text{lpo}} h(t)$.

Assignment 5 (Knuth–Bendix completion)

(14 points)

Let E be the following set of equations over $\Sigma = (\{f/1, g/1, h/1, b/0\}, \emptyset)$.

$$f(g(x)) \approx h(x) \quad (1)$$

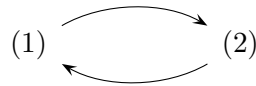
$$g(f(x)) \approx h(x) \quad (2)$$

Apply the Knuth-Bendix completion procedure to E and transform it into a finite convergent term rewrite system; use a Knuth-Bendix ordering with weight 1 for all function symbols and variables and the precedence $f > g > h > b$. Use a reasonable strategy.

Assignment 6 (Dependency pairs)

(4 + 4 + 4 = 12 points)

Part (a) Find an example of a terminating rewrite system R with exactly two rewrite rules such that $\text{DP}(R)$ contains exactly two dependency pairs and the dependency graph of R has the shape



Specify both R and $\text{DP}(R)$.

Part (b) Find an example of a terminating rewrite system R with exactly two rewrite rules such that $\text{DP}(R)$ contains exactly two dependency pairs and the dependency graph of R has the shape



Specify both R and $\text{DP}(R)$.

Part (c) Find an example of a terminating rewrite system R with exactly two rewrite rules such that $\text{DP}(R)$ contains exactly two dependency pairs and the dependency graph of R has the shape



Specify both R and $\text{DP}(R)$.