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Tutorials for “Logic in Computer Science”  
 Exercise sheet 9

**Exercise 9.1:**

If a list  $l$  has the form  $[a_1, \dots, a_n]$ , then  $a_{i+1}$  is called in successor of  $a_i$  in  $l$ . (For instance, **a**, **r**, and **n** are successors of **a** in the list  $[s, a, a, r, l, a, n, d]$ .) Implement a Prolog predicate `succ( $l, x, y$ )` that succeeds if  $y$  is a successor of  $x$  in  $l$ . Can your implementation also be used to compute the predecessors of an element in a list?

**Notes on Prolog programming:** You can find several Prolog implementations on the computers at the University campus. For instance, at the CIP pool in the first floor of Bldg. 45 you can find SWI-Prolog (`/usr/local/bin/pl`); on the MPI computers, GNU Prolog is available (`/opt/gnu/bin/gprolog`). Sources and binaries (Linux, MS-Windows, MacOS X) for SWI-Prolog can also be downloaded from <http://www.swi-prolog.org/>.

To use Prolog, write your clauses into a file `filename.pl`, start Prolog, and type

`[ 'filename.pl' ] .`

(including the square brackets, the quotes, and the period at the end). Queries can be entered directly at the Prolog prompt (also terminated with a period). If a Prolog system has found a solution (i.e., a success node of the SLD-tree), you can type a semicolon to start the search for another solution.

A remark on the Prolog syntax: To prevent typos, most Prolog implementations issue a warning message if a program clause contains a *singleton variable* (i.e., a variable that occurs only once). To avoid this, use an underscore (`_`) instead of a regular variable name for singleton variables.

**Exercise 9.2:**

Define a *rotation* of a list  $l$  as follows:

- If  $l$  is the empty list  $[]$ , then  $l$  is a rotation of  $l$ .
- If  $l$  is a list  $[a_1, \dots, a_n]$  with  $n \geq 1$ , then  $l$  is a rotation of  $l$ , and every list  $[a_i, \dots, a_n, a_1, \dots, a_{i-1}]$  with  $1 < i \leq n$  is a rotation of  $l$ .

The following Prolog program computes rotations of a list. (We use a predicate `append1` rather than `append`, since the latter is predefined in most Prolog systems.)

```
append1([],L2,L2).
append1([X|L1],L2,[X|L3]) :- append(L1,L2,L3).

rotate(L,L).
rotate([X|L],R) :- append1(L,[X],R1), rotate(R1,R).
```

Describe the SLD-tree for the query `rotate([a,b],X)` and show that it is infinite.

**Exercise 9.3:**

Give an alternative implementation for `rotate` so that, given any ground list  $l$ , the SLD-tree for the query `rotate(l,X)` is finite.

Hint: `append1` can not only be used to append two lists, but also to split a list in two sublists.

Challenge: Implement `rotate` in such a way that every rotation of a list is computed exactly once (i.e., if  $l$  is a ground list of length  $n$ , then the SLD-tree for `rotate(l,X)` is finite and contains exactly one success node for  $n = 0$  and exactly  $n$  success nodes for  $n > 0$ .)

**Exercise 9.4:**

Is it possible that a Prolog system terminates for a query  $G_1$ , terminates for a query  $G_2$ , and loops for the query  $G_1, G_2$ ? If so, how?

Put your solution into the mail box at the door of room 627 in the MPI building (46.1) before June 21, 11:00 (Group D: before June 24, 11:00). Don't forget to write your name and the name of your tutorial group (B, C, D) on your solution.