

**Assignment 1 (OBDDs)**

(3 + 6 + 3 = 12 points)

Let  $F$  be the propositional formula  $P \wedge (Q \vee R) \wedge S$ .

**Part (a)** Give the reduced OBDD for  $F$  w. r. t. the ordering  $P < Q < R < S$ .

**Part (b)** Find a total ordering over  $\{P, Q, R, S\}$  such that the reduced OBDD for  $F$  has 6 non-leaf nodes. Give the resulting reduced OBDD.

**Part (c)** For how many total orderings over  $\{P, Q, R, S\}$  does the reduced OBDD for  $F$  have 6 non-leaf nodes?

**Assignment 2 (Ordered Resolution)**

(7 + 7 = 14 points)

In Sect. 3.13 of the lecture notes (p. 79), the inference rules for ground resolution with ordering restrictions (without selection functions) are given by

(Ground) Ordered Resolution:

$$\frac{D \vee A \quad C \vee \neg A}{D \vee C} \quad \text{if } A \succ L \text{ for all } L \text{ in } D \text{ and } \neg A \succeq L \text{ for all } L \text{ in } C.$$

(Ground) Ordered Factorization:

$$\frac{C \vee A \vee A}{C \vee A} \quad \text{if } A \succeq L \text{ for all } L \text{ in } C.$$

This calculus is sound and refutationally complete for sets of ground clauses.

Suppose that we replace the ordering restriction for the first inference rule by “if  $A \succ L$  for all  $L$  in  $D$  and  $A \succeq L$  for all  $L$  in  $C$ .”

**Part (a)** Is the calculus with this modification still sound? If yes, give a short explanation; if no, give a counter example.

**Part (b)** Is the calculus with this modification still refutationally complete? If yes, give a short explanation; if no, give a counter example.

**Assignment 3 (E-Algebras)**

(12 points)

Let  $\Sigma = (\Omega, \emptyset)$  be a first-order signature with  $\Omega = \{f/1, b/0, c/0, d/0\}$ . Let  $E$  be the set of  $\Sigma$ -equations

$$\{\forall x (f(x) \approx b), c \approx d\},$$

let  $X = \{x, y, z\}$  be a set of variables. For any  $t \in T_\Sigma(X)$  let  $[t]$  denote the congruence class of  $t$  w.r.t.  $E$ . Let  $\mathcal{T} = T_\Sigma(X)/E$ , let  $U_\mathcal{T}$  be the universe of  $\mathcal{T}$ , and let  $\beta : X \rightarrow U_\mathcal{T}$  be the assignment that maps every variable to  $[c]$ . Decide for each of the following statements whether they are true or false:

- |  |  |
|--|--|
| (1) $[z]$ is a finite set of $\Sigma$ -terms.    | (5) $U_\mathcal{T}$ is finite.                             |
| (2) $[f(z)]$ is a finite set of $\Sigma$ -terms. | (6) $[b] \in U_\mathcal{T}$ .                              |
| (3) $[c]$ is a set of ground $\Sigma$ -terms.    | (7) $\{x, y\} \in U_\mathcal{T}$ .                         |
| (4) $[f(c)]$ is a set of ground $\Sigma$ -terms. | (8) $\mathcal{T}(\beta)(\forall z (z \approx f(x))) = 1$ . |

(Note on grading: A yes/no answer is sufficient; you do not have to give any explanations. However, you need at least five correct answers to get any points for assignment 3. Missing answers count like false answers.)

**Assignment 4** (*Term Orderings*) (7 + 7 = 14 points)

Let  $\Sigma = (\Omega, \Pi)$  be a finite signature, let  $\succ$  be a strict partial ordering on  $\Omega$ , let  $s, t \in T_\Sigma(X)$ .

**Part (a)** Prove: If  $s$  contains a subterm  $s' = f(s_1, \dots, s_n)$  such that  $\text{var}(s') \supseteq \text{var}(t)$  and  $f \succ g$  for all function symbols  $g$  occurring in  $t$ , then  $s \succ_{\text{lpo}} t$ .

**Part (b)** Refute: If  $s$  contains a subterm  $s' = f(s_1, \dots, s_n)$  such that  $\text{var}(s) \supseteq \text{var}(t)$  and  $f \succ g$  for all function symbols  $g$  occurring in  $t$ , then  $s \succ_{\text{lpo}} t$ .

**Assignment 5** (*Critical Pairs, Termination*) (8 + 6 = 14 points)

Let  $E$  be the following set of equations over  $\Sigma = (\{f/1, g/2, h/2, b/0, c/0\}, \emptyset)$ .

$$f(f(x)) \approx g(b, x) \quad (1)$$

$$h(f(y), y') \approx f(h(y), y') \quad (2)$$

$$g(h(z, z), c) \approx h(z, b) \quad (3)$$

**Part (a)** Suppose that the three equations in  $E$  are turned into rewrite rules by orienting them from left to right. Give all (non-trivial) critical pairs between the resulting three rules.

**Part (b)** It is possible to orient the equations in  $E$  using an appropriate KBO, so that there are no critical pairs between the resulting rules. Give the weights and precedence of the KBO; explain how the equations are oriented.

**Assignment 6** (*Dependency Pairs*) (3 + 8 + 3 = 14 points)

**Part (a)** Let  $\Sigma = (\Omega, \emptyset)$  with  $\Omega = \{f/2, g/2, h/1, k/1, b/0\}$ . Compute the dependency pairs of the following rewrite system  $R$  over  $\Sigma$ :

$$f(x, y) \rightarrow g(k(y), y) \quad (1)$$

$$g(h(x), y) \rightarrow f(x, y) \quad (2)$$

$$g(x, h(y)) \rightarrow h(f(x, y)) \quad (3)$$

$$k(x) \rightarrow x \quad (4)$$

$$k(h(x)) \rightarrow k(x) \quad (5)$$

**Part (b)** Compute the approximated dependency graph for  $R$  (using cap and ren) and use the subterm criterion to reduce the graph as much as possible. If a graph is modified, depict both the original and the modified graph and indicate the strongly connected components in the graphs.

**Part (c)** There is one cycle  $K$  in the approximated dependency graph for  $R$  that cannot be eliminated using the subterm criterion. In fact, there is a  $K$ -minimal infinite rewrite sequence for this cycle. How does it look like?