

**Assignment 1** (*Propositional Logic*)

(10 points)

Let  $F[G \wedge H]$  be a propositional formula that contains  $G \wedge H$  as a subformula (where  $G$  and  $H$  are also propositional formulas). Prove: If  $F[G \wedge H]$  is valid, then  $G \rightarrow F[H]$  is valid.

**Assignment 2** (*Resolution*)

(8 + 8 = 16 points)

Let  $\Sigma = (\{a/0, b/0, f/1, g/1\}, \{P/2, Q/1, R/1, S/1\})$ ; let  $N$  be the following set of clauses over  $\Sigma$ :

$$\neg Q(y) \vee S(x) \vee P(x, x) \vee P(y, g(y)) \quad (1)$$

$$\neg P(z, g(a)) \vee R(z) \quad (2)$$

$$\neg S(a) \vee \neg S(f(b)) \quad (3)$$

$$S(f(y)) \vee S(y) \quad (4)$$

**Part (a)**

Suppose that the atom ordering  $\succ$  is an LPO with the precedence  $P > Q > R > S > f > g > a > b$ . Compute all ordered resolution inferences between the clauses (1)–(4) with respect to  $\succ$ . (Compute only inferences between the clauses given here, not between derived clauses. Do not compute any inferences that violate the ordering conditions of ordered resolution.)

**Part (b)**

If a selection function is defined appropriately, the set  $N$  is saturated under ordered resolution with selection (w.r.t. the ordering  $\succ$  from Part (a)). Which literals have to be selected?

**Assignment 3** (*Tableaux*)

(10 points)

Use semantic tableaux to show that the following set of formulas over  $\Sigma = (\{b/0, f/1\}, \{P/2\})$  is unsatisfiable:

$$\begin{aligned} \forall y \forall x \left( P(x, y) \rightarrow P(f(x), f(y)) \right) \\ \exists w P(b, w) \\ \forall z \neg P(f(f(b)), z) \end{aligned}$$

Use exactly the expansion rules given in the lecture; do not use shortcuts.

**Assignment 4** (*E-Algebras*)

(6 + 8 + 6 = 20 points)

Let  $\Sigma = (\Omega, \emptyset)$  with  $\Omega = \{a/0, b/0, f/1\}$ ; let  $E$  be the set of equations  $\{a \approx b, f(b) \approx f(f(b))\}$ .

**Part (a)**

Show that  $f(a) \leftrightarrow_E^* f(f(f(a)))$ .

**Part (b)**

How many elements does the universe of the quotient algebra  $T_\Sigma(\emptyset)/E$  have?

**Part (c)**

Give an example of a (quantified) equation  $\forall \vec{x} (t \approx t')$  such that  $t \leftrightarrow_E^* t'$  does not hold, but  $T_\Sigma(\emptyset)/E \models \forall \vec{x} (t \approx t')$ .

**Assignment 5** (*Reduction Orderings*)

(8 + 8 = 16 points)

For a signature  $\Sigma$  we define  $T_x^1$  as the set of all  $\Sigma$ -terms that contain exactly one occurrence of the variable  $x$  and no other variables.

**Part (a)**

Prove: If all function symbols in  $\Sigma$  have arity 1, then a Knuth-Bendix ordering  $\succ$  with a total precedence is total on  $T_x^1$ .

**Part (b)**

Prove: If  $\Sigma$  contains a binary function symbol and a constant function symbol, then there exists no reduction ordering that is total on  $T_x^1$ .

**Assignment 6** (*Feature Vector Indexing*)

(8 points)

Decide for each of the following numbers whether or not it could be used as a feature in a feature vector index:

- (1) the number of ground arguments of predicate symbols in a clause,
- (2) the number of variable occurrences in a clause,
- (3) the number of constant symbols occurring in positive literals in a clause,
- (4) the absolute value of the difference between the number of positive and the number of negative literals in a clause.

(Note on grading: A yes/no answer is sufficient; you do not have to give an explanation. However, you need at least three correct answers in this assignment to get any points. Missing answers count like false answers.)