

**Problem 1** (*Semantics*)

(14 points)

Show that the following inference rule is sound:

$$\frac{D \vee f(s) \approx s' \quad C \vee f(t) \approx t'}{D \vee C \vee s \not\approx t \vee s' \approx t'}$$

where the premises and conclusions are (implicitly universally quantified) equational clauses and the two premises have no common variables.

**Problem 2** (*Ordered resolution with selection*)

(10 points)

Let  $\Sigma = (\Omega, \Pi)$  be a signature with  $\Omega = \{a/0, b/0, f/1, g/1\}$  and  $\Pi = \{p/1, q/1, r/2\}$ . Let  $N$  be the following set of clauses:

$$\neg p(g(a)) \tag{1}$$

$$q(f(x)) \tag{2}$$

$$r(x, x) \tag{3}$$

$$\neg q(g(x)) \vee p(g(x)) \tag{4}$$

$$\neg r(a, z) \vee q(a) \vee p(f(z)) \tag{5}$$

$$\neg r(f(y), y) \vee \neg q(y) \vee \neg p(f(f(y))) \tag{6}$$

Suppose that an ordering  $\succ$  on ground atoms is defined in such a way that  $p(\dots) \succ q(\dots) \succ r(\dots)$ . Which literals in which clauses must be selected by a selection function  $S$  such that  $N$  is saturated under  $Res_{\Sigma}^{\succ}$ ?

**Problem 3** (*Tableaux*)

(16 points)

Use semantic tableaux to show that the following set of formulas is unsatisfiable:

$$\begin{aligned} &\forall x \left( \exists y p(x, y) \rightarrow \exists z p(f(x), z) \right) \\ &\quad p(a, a) \\ &\quad \neg \exists x p(f(f(a)), x) \end{aligned}$$

(Note: Quantifiers extend over the shortest following subformula, or in other words,  $\exists y F \rightarrow \exists z G$  means  $(\exists y F) \rightarrow (\exists z G)$ .)

**Problem 4** (*E-algebras*)

(8 + 8 = 16 points)

Let  $\Sigma = (\Omega, \emptyset)$  with  $\Omega = \{a/0, b/0, f/1\}$ ; let  $E$  be the set of (implicitly universally quantified) equations  $\{f(f(x)) \approx a\}$ .

**Part (a)**

Give one possible derivation for the statement  $E \vdash f(a) \approx a$ .

**Part (b)**

Is the universe of the initial  $E$ -algebra  $T_\Sigma(\emptyset)/E$  finite or infinite? If it is finite, how many elements does it have?

**Problem 5** (*Reduction orderings*)

(10 points)

Let  $\Sigma$  be an arbitrary first-order signature. Define the ordering  $\succ$  on  $T_\Sigma(X)$  by

$$s \succ s' \text{ if and only if } |s| > |s'|$$

where  $|t|$  is the size of  $t$ , that is, the cardinality of  $\text{pos}(t)$ . Is  $\succ$  a reduction ordering? Give a proof or a counterexample.

**Problem 6** (*Knuth-Bendix completion*)

(14 points)

Let  $\Sigma = (\Omega, \emptyset)$  with  $\Omega = \{a/0, b/0, f/1, g/2, h/1\}$ , let  $\succ$  be the LPO with precedence  $h > g > f > b > a$ , and let  $E$  be the following set of equations:

$$f(x) \approx g(h(x), x) \quad (1)$$

$$g(a, b) \approx b \quad (2)$$

$$h(b) \approx a \quad (3)$$

Use Knuth-Bendix completion to transform  $E$  into an equivalent convergent set  $R$  of rewrite rules such that  $\rightarrow_R \subseteq \succ$ .